HW6

1. Consider the Conjugate Gradient algorithm in class. The search directions are obtained as follows

$$p^{(0)} = r^{(0)} = b - Ax^{(0)}.$$

For later directions, one construct

$$p^{(k+1)} = r^{(k+1)} + s_k p^{(k)}$$

One chooses s_k such that $p^{(k+1)}$ is A-orthogonal to $p^{(k)}$. Fill in the details (and repeat) in the proof in lecture notes to

- Show that p^i 's are A-orthogonal search directions.
- Show that the formula for s can be simplified to the following.

$$s_k = \frac{|r^{(k+1)}|^2}{|r^{(k)}|^2}.$$

Find the lecture notes for the proof.

2. Find the singular value decomposition of

$$A = \left(\begin{array}{rrrr} 1 & 0 & 1 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & 0 \end{array}\right)$$

- 3. Prove the following using singular value decomposition
 - Show that the 2-norm of A is the largest singular value and the Frobenius norm is the square root of the sum of squares of all singular values.

Consider the SVD:

$$A = U\Sigma V^T$$

where U is an $m \times m$ orthogonal matrix, V is an $n \times n$ orthogonal matrix, and $\Sigma_{ij} = \sigma_i \delta_{ij}$ with $\sigma_1 \ge \sigma_2 \ge \cdots \sigma_{\min(m,n)} \ge 0$. Let $D = \Sigma^T \Sigma = (\sigma_i^2 \delta_{ij})_{n \times n}$ and it is thus an $n \times n$ square matrix. We have assumed $\sigma_i = 0$ for $i \ge \min(m, n)$. The 2-norm is given

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by

$$\|A\|_{2} = \sqrt{\rho(A^{T}A)} = \sqrt{\rho(V\Sigma^{T}U^{T}U\Sigma V^{T})} = \sqrt{\rho(VDV^{T})}$$
$$= \sqrt{\rho(D)} = \sigma_{1}$$

Above, we have used the fact $\lambda(VDV^T) = \lambda(D)$. For the Frobenius norm, we have

$$\begin{split} \|A\|_F &= \sqrt{\sum_{i,j} a_{ij}^2} = \sqrt{\operatorname{tr}(A^T A)} = \sqrt{\operatorname{tr}(V D V^T)} \\ &= \sqrt{\operatorname{tr}(D V^T V)} = \sqrt{\operatorname{tr}(D)} = \sqrt{\sum_i \sigma_i^2}. \end{split}$$

 Show that AA^T and A^TA have the same 2 norm. Note that AA^T is not the transpose of A^TA! Using the same SVD as above, we have

$$AA^T = U\Sigma\Sigma^T U^T, \quad A^T A = VDV^T.$$

Since AA^T is real symmetric, the 2-norms is the largest eigenvalue, or the spectral radius, given by $\rho(\Sigma\Sigma^T) = \sigma_1^2$. Similarly, $||A^TA||_2 = \rho(D) = \sigma_1^2$.

4. Write a code for power method and apply it to the following matrix to find the leading eigen-pair:

$$\left(\begin{array}{rrrr} -4 & 14 & 0\\ -5 & 13 & 0\\ -1 & 0 & 2 \end{array}\right)$$

5. Suppose matrix A is a real matrix with eigenvalues $|\lambda_1| > |\lambda_2| \ge \cdots \ge |\lambda_n|$. Show that the power method still works even if the matrix A does not have n eigenvectors (you may consider the Jordan canonical form).

According to power method,

$$x^{(k)} = \frac{1}{\mu_k} A^k x_0,$$

where μ_k is some scalar to make the $\|\cdot\|_{\infty}$ norm to be 1. Clearly, the point is to check how $A^k x_0$ behaves. We want to say that the component corresponding of λ_1 dominates.

We use the Jordan canonical form:

$$A = PJP^{-1}.$$

According to the assumption, we know

$$J = \begin{pmatrix} \lambda_1 & & & \\ & J_{\mu_2} & & \\ & & \ddots & \\ & & & J_{\mu_m} \end{pmatrix}$$

where J_{μ_k} are some Jordan blocks. Hence, the first column of P is the eigenvector corresponding to λ_1 . Though there are no enough eigenvectors, we use the columns of P, $\{v_1, \dots, v_n\}$, as the basis (the generalization of eigenvectors):

$$A^k x_0 = P J^k P^{-1} \left(\sum_{i=1}^n \alpha_i v_i\right)$$

Since $I = P^{-1}P = P^{-1}[v_1, \dots, v_n]$, we find $P^{-1}v_i = e_i$ (the vector with *i*th component being 1 and others being zero.) Hence,

$$A^{k}x_{0} = \sum_{i=1}^{n} \alpha_{i}(PJ^{k})_{i} = \alpha_{1}\lambda_{1}^{k}v_{1} + \sum_{i=2}^{n} \alpha_{i}w_{i,k}$$

where B_i means the *i*th column of B, and $w_{i,k} = (PJ^k)_i = P(J^k)_i$. The *i*th column of J^k corresponds to one Jordan block: $J^k_{\mu_\ell}$ for some ℓ . Direct computation shows

$$J_{\mu_{\ell}}^{k} = (\mu_{\ell}I + F)^{k} = \sum_{j=0}^{\min(n,k)} C_{k}^{j} \mu_{\ell}^{k-j} F^{j}.$$

This shows that

$$|w_{ik}| \le C \sum_{j=0}^{n} |k|^{j} |\mu_{\ell}|^{k-j} \ll |\lambda_{1}|^{k}.$$

Note that μ_{ℓ} is λ_i . In other words,

$$\frac{|A^k x_0 - \alpha_1 \lambda_1^k v_1|}{|\lambda_1|^k} \to 0.$$

The eigenvector then converges. Hence, the power method works.

6. This is about Householder transform

(a) The Householder transform

$$P = I - 2w^T w$$

is the identity matrix perturbed a rank 1 matrix. Show that n-1 eigenvalues are unchanged while one eigenvalue is changed to be -1. (Hint: use the geometric meaning of the transform. Try to identify the eigenvectors using the geometric meaning directly) For identity matrix, all eigenvalues are 1. With the rank 1 perturbation $-2ww^{T}$, one eigenvalue is changed. To figure out how this is changed, we check the geometric meaning.

According to the geometric meaning, w is an eigenvector with egienvalue -1, while all the nonzero vectors perpendicular to w are eigenvectors with eigenvalue 1. Since the space perpendicular to w has dimension n - 1, we thus can find n - 1 independent such eigenvectors. Hence, all the n eigenvectors are found.

(b) Let H_1 be a Householder transform in \mathbb{R}^{n-1} . Show that transform

$$Q_1 = \left(\begin{array}{cc} 1 & 0\\ 0 & H_1 \end{array}\right)$$

is also a Householder transform.

Let $H_1 = I_{n-1} - 2w_1w_1^T$ for some $w_1 \in \mathbb{R}^{n-1}$ with $|w_1| = 1$. According to the formula, the matrix Q_1 is doing the following: the first component of x is unchanged, while for the other n-1components, it is doing the reflection by H_1 . Geometrically, it should be a reflection in \mathbb{R}^n as well. This motivates us to consider

$$w = \left(\begin{array}{c} 0\\ w_1 \end{array}\right).$$

Hence, we guess:

$$Q_1 = I_n - 2ww^T.$$

In fact, one can check this easily:

$$I_n - 2ww^T = \begin{pmatrix} 1 & 0 \\ 0 & I_{n-1} \end{pmatrix} - 2 \begin{pmatrix} 0 & 0 \\ 0 & w_1 w_1^T \end{pmatrix} = Q_1.$$

Lastly, it is easy to see $|w| = |w_1| = 1$. The claim is then proved.

7. Give a 2×2 matrix such that the diagonal elements of A_k in the QR algorithm do not converge to the eigenvalues.

You can construct a matrix with complex eigenvalues. Other examples are also possible, for example, many of you constructed A = [0, 1; 1, 0]

8. Apply QR algorithm (you may code) to find all eigenvalues of the following matrix (with accuracy $10^{-4})$

$$\left(\begin{array}{rrrr} 3 & 1 & 0 \\ 1 & 4 & 2 \\ 0 & 2 & 3 \end{array}\right)$$